# The geometry of synchronisation problems and group actions

Jacek Brodzki Centre for Geometry, Topology, and Applications University of Southampton Análisis Topológico de Datos November 2018

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- Synchronisation in examples
- Spectral properties of graphs
- Spectral clustering of graphs
- The geometry and topology of synchronisation
- Learning group actions
- Persistence

## Synchronisation in examples

## Synchronisation

Synchronisation problems deal with questions regarding aligning a collection of objects in a consistent manner.

Given a sequence of objects  $(o_1, ..., o_n)$  and a group (or collection of groups) *G* the objective is to learn a collection of group elements  $\rho_{ij}$  that would transform object  $o_i$  to  $o_j$  in a consistent way.

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We'd like to get as close as possible to a group action, so if  $\rho_{ij}: o_i \mapsto o_j$  and  $\rho_{jk}: o_j \mapsto o_k$  then

 $\rho_{ik} = \rho_{ij} \circ \rho_{jk} : o_i \mapsto o_k.$ 

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If it is possible to find a sequence of group elements that allow for an accurate transformation between objects then this set can be *synchronised*.

#### **Examples of synchronisation**

- Orientability of manifolds [Singer, Wu (2011)]: G = O(1);
- Angular synchronisation [Singer (2011)]: G = U(1);
- Vector diffusion maps [Singer, Wu (2012)]: G = O(d);
- Collection Shape Matching [Nguyen et al. (2011)], [Huang, Guibas (2013)], [Chen et al. (2014)], [Maron et al. (2016)]:
  G = E<sup>d</sup>;
- Cryo-EM Structural Reconstruction [Singer et al. (2011)], [Shkolnisky, Singer (2012)], [Zhao, Singer (2014)], [Bandeira et al. (2015)]: G = SO(3)

A differentiable manifold M is a topological space equipped with an atlas  $\{U_{\alpha}\}$  which is an open cover of M such that for every  $\alpha$ there is a homeomorphism

 $\phi_{\alpha}: U_{\alpha} \to \mathbb{R}^n$ 

which are such that for every nonempty overlap  $U_{lpha} \cap U_{eta}$  the transition map

$$g_{lphaeta}=\phi_lpha\circ\phi_eta^{-1}:\phi_eta(U_lpha\cup U_eta) o\phi_lpha(U_lpha\cup U_eta)$$

is a diffeomorphism of open subsets of  $\mathbb{R}^n$ .

A manifold is orientable iff it admits an oriented atlas, i.e., an atlas  $\{U_{\alpha}\}$  such that all the transition maps  $g_{\alpha\beta}$  are orientable.

This in turn means that the Jacobian of each such map has an everywhere positive determinant.

#### **Question:**

How can one determine if a given manifold is orientable starting from a sample of points of M?

Assume given points  $S = \{x_1, x_2, ..., x_n\}$  sampled from some  $\mathbb{R}^p$ ; we assume that these points lie on a submanifold M of  $\mathbb{R}^p$ .



What is a good notion of tangent space in this case?

#### Principal Component Analysis PCA

- PCA was invented by Karl Pearson in 1901
- It provides a way of reducing the dimension of data, which is particularly useful when trying to visualise data in two or three dimensions.



Choose  $x_i \in S$  and identify its K nearest neighbours, where K << n. These points  $x_{i_1}, \ldots, x_{i_K}$  are located near the tangent space  $T_{x_i}M$ . There may of course be deviations due to noise or curvature of M.



#### Local PCA

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Let  $\mu$  be the average of the nearest neighbours of  $x_i$ :

$$u = \frac{1}{K} \sum_{k=1}^{K} x_{i_k}$$

and form the matrix of vectors

$$X_i = [x_{i_1} - \mu \quad x_{i_2} - \mu \quad \dots \quad x_{i_K} - \mu].$$

So the *k*-th column is the vector  $x_{i_k} - \mu$  and the matrix is of size  $p \times K$ .

$$X_i = [x_{i_1} - \mu \quad x_{i_2} - \mu \quad \dots \quad x_{i_K} - \mu].$$

Assume that the dimension of M from which the points  $x_i$  are sampled is d.

If the nearest points  $x_{i_k}$  to  $x_i$  are elements of  $T_{x_i}M$  (more precisely: the vectors  $x_{i_k} - x_i$  are elements of  $T_{x_i}M$ ) then the matrix X will have d linearly independent columns, i.e., its rank is d.

This can be detected through the singular values of  $X_i$ .

Assume A is a rectangular  $m \times n$  matrix;  $A \in \mathbb{C}^{m \times n}$ , with  $m \ge n$ .

A singular value decomposition of A is

$$A = U \begin{pmatrix} \Sigma \\ 0 \end{pmatrix} V^*$$

where

$$\Sigma = \begin{pmatrix} \sigma_1 & & \\ & \ddots & \\ & & \sigma_n \end{pmatrix}, \quad \sigma_1 \ge \sigma_2 \ge \cdots \ge \sigma_n \ge 0,$$

and  $U \in \mathbb{C}^{m \times m}$  and  $V \in \mathbb{C}^{n \times n}$  are unitary. The numbers  $\sigma_k$  are called *singular values*.

The matrix U is a left singular vector matrix, and V is the right singular vector matrix.

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 is an SVD decomposition of A then

 $A^* = V(\Sigma \quad 0)U^*$ 

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- A is nonsingular iff all its singular values are nonzero.
- The rank of a matrix is the number of its nonzero singular values.
- Vectors of *U* corresponding to the nonzero singular values span the column space of *A* (the image of *A* regarded as a linear operator).

$$\begin{pmatrix} 1 & 2 \\ 3 & 4 \\ 5 & 6 \\ 7 & 8 \end{pmatrix} = \begin{pmatrix} -0.1525 & -0.8226 & -0.3945 & -0.3800 \\ -0.3499 & -0.4214 & 0.2428 & 0.8007 \\ -0.5474 & -0.0201 & 0.6979 & -0.4614 \\ -0.7448 & 0.3812 & -0.5462 & 0.0407 \end{pmatrix} \\ \times \begin{pmatrix} 14.2691 & 0 \\ 0 & 0.6268 \\ 0 & 0 \\ 0 & 0 \end{pmatrix} \\ \times \begin{pmatrix} -0.6414 & 0.7672 \\ -0.7672 & -0.6414 \end{pmatrix}$$

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Choose a threshold  $\gamma \in [0, 1]$ . In practice  $\gamma$  is closer to 1 than 0. We estimate the dimension  $d_i$  of  $T_{x_i}M$  as the smallest integer  $d_i$  for which

$$\frac{\sum_{k=1}^{d_i} \sigma_{i,k}^2}{\sum_{k=1}^K \sigma_{i,k}^2} > \gamma.$$

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E.g., if  $\gamma = 0.9$  then  $d_i$  singular values account for at least 90% of the variability of the data, but  $d_i - 1$  singular values will account for less than 90%.

Having estimated the local dimension  $d_i$  for all i = 1, ..., n we need an estimate of the global dimension. Possible options are:

The mean

$$d = \frac{1}{n} \sum_{i=1}^{n} d_i$$

rounding to the nearest integer. This minimises the sum of squared errors

$$\sum_{i=1}^n (d_i-d)^2.$$

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#### The median

$$d =$$
median  $(d_1, d_2, \ldots, d_n).$ 

This minimises the  $\ell^1\text{-}\text{error}$ 

$$\sum_{i=1}^n |d_i - d|.$$

This is more robust with respect to outliers than the mean estimator.

Consider again the matrix describing the shifted local neighbourhood of  $x_i$ :

$$X_i = [x_{i_1} - \mu \quad x_{i_2} - \mu \quad \dots \quad x_{i_K} - \mu].$$

Assume that the matrix  $X_i$  has an SVD decomposition

$$X_i = U_i \Sigma_i V_i^t,$$

where the columns of the  $p \times K$  matrix  $U_i$  (the left singular vectors) are orthonormal.

#### Local coordinate frame

Define the matrix  $O_i$  by

$$O_i = [u_{i_1}u_{i_2}\ldots u_{i_d}]$$

where the column vectors  $u_{i_j}$  are the first d left singular vectors of the matrix  $X_i$  (arranged in the decreasing order of the singular values)

The columns of  $O_i$  provide an approximation of an orthonormal basis of the tangent space  $T_{x_i}(M)$ .



Now consider two nearby points  $x_i$  and  $x_j$  which which are each other's nearest neighbours. We have identified their local frames, can we map one to the other?



#### Alignment

We have found matrices  $O_i$  and  $O_j$  approximating the tangent spaces at each point. Because of possible noise or curvature effects, the matrix  $O_i^T O_j$  need not be orthogonal, and we define  $O_{ij}$  to be an orthogonal matrix which is closest to  $O_i^T O_j$ :

$$O_{ij} = \operatorname{argmin} \| O - O_i^T O_j \|$$

where the minimum is taken over all orthogonal matrices  $O \in O(d)$ .

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This has the following solution.

If the SVD of  $O_i^T O_j$  is

$$O_i^T O_j = U \Sigma V^T$$

then

$$O_{ij} = UV^T$$
.

The process of finding optimal orthogonal transformations between bases is called *alignment*.

Given that  $O_{ij} \in O(d)$ , we have that det  $O_{ij} = \pm 1$  where

- det  $O_{ij} = 1$  iff  $O_{ij} \in SO(d)$ , so it is a 'rotation'
- det  $O_{ij} = -1$  iff  $O_{ij}$  is a composition of a rotation and a reflection.

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Note that this only works between nearby points, and the outcome can be encoded as a graph  $\mathcal{G} = (V, E)$  with vertices corresponding to the *n* data points, and where an edge (i, j) exists iff there is a corresponding matrix  $O_{ij}$ .
### Synchronisation

We can encode the results of our investigation so far by defining a matrix

$$Z_{ij} = egin{cases} { ext{det}} \ O_{ij}, & ext{if} \ (i,j) \in E \ 0, & ext{otherwise} \end{cases}$$

For each point  $x_i$  we have a basis defined by local *PCA* as well as a basis given by the ambient manifold.

Define:

$$z_i = egin{cases} 1, & ext{if the orientations agree} \ -1, & ext{otherwise.} \end{cases}$$

Clearly,

$$z_i z_j^{-1} = Z_{ij}.$$

In practice, because of the curvature and noise, the relation

$$z_i z_j^{-1} = Z_{ij}.$$

will not necessarily hold for all edges.

Synchronisation in this case refers to finding the function z on the vertices (z takes values in  $\mathbb{Z}_2$ ) for which this equation holds. When the manifold is orientable, this solution exists, while for nonorientable manifolds we will find two classes of frames connected by matrices of determinant 1.

#### Angular synchronisation

#### Problem

Find an accurate estimation of a set of unknown angles  $\theta_1, \ldots, \theta_n$ from *m* noisy measurements of the differences  $\theta_i - \theta_i$ .

A main motivation for this problem comes from systems of coupled oscillators. This is relevant, e.g., in power grids where generators on the network need to be synchronised very precisely.



Other applications include

- Time synchronisation of distributed networks given the differences t<sub>i</sub> − t<sub>j</sub> from which we need to determine the times t<sub>1</sub>,..., t<sub>n</sub> ∈ ℝ.
- Surface reconstruction in computer vision and optics, where the surface is reconstructed from noisy measurements of the gradient of the surface.

#### Angular synchronisation

Main problem: given the set of differences  $\theta_i - \theta_j$  (call them offsets), we do not know which of them are reliable, and which are noise.

So there are two kinds of edges in the synchronisation graph  $\mathcal{G}$ :

$$\begin{split} \delta_{ij} &= \theta_i - \theta_j & (i,j) \in E_{\text{good}} \\ \delta_{ij} &\sim \text{Uniform}([0,2\pi)) & (i,j) \in E_{\text{bad}}. \end{split}$$

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We note that cycles in the graph G satisfy the following consistency relation. For a triangle of 'good' edges (i, j), (j, k), and (k, i) we have

$$\delta_{ij} + \delta_{jk} + \delta_{ki} = 0 \mod 2\pi.$$

This will be quite important in later lectures.

The following new framework emerges from our examples:

- Similarity between data points is described by a graph of relationships *G*;
- If two points are similar enough, they are connected by and edge with a weight that quantifies their affinity.
- Many examples show that apart from a weight, it may be useful to have an orthogonal transformation attached to an edge.



To each pair we would like to assign an orthogonal transformation  $O_{ij}$  that best aligns the images and a weight  $w_{ij}$  describing their measure of affinity.

Given a graph  $\mathcal{G} = (V, E)$  of relationships define

$$\mathcal{S}_{i,j} = egin{cases} w_{ij} O_{ij}, & (i,j) \in E \ 0 & ext{otherwise} \end{cases}$$

where  $O_{ij} \in O(d)$ ,  $0 \in O(d)$ . We assume  $w_{ij} = w_{ji}$  and  $O_{ji} = O_{ij}^T$ .

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where  $O_{ij} \in O(d)$ ,  $0 \in O(d)$ . We assume  $w_{ij} = w_{ji}$  and  $O_{ji} = O_{ij}^T$ . As previously, we assume that we have identified tangent spaces at data points in V and want to create a suitable 'parallel transport' operator to compare the tangent spaces at nearby points. These are called *diffusion maps*. Vertex degree:

$$deg(i) = \sum_{j \sim i} w_{ij}.$$
$$D(i, i) = deg(i)I_{d \times d}$$

Now for  $x_i \in V$  consider a vector  $v(i) \in \mathbb{R}^d$  regarded as an element of  $\mathbb{R}^d$ .

'Divergence'

$$(D^{-1}Sv)(i) = \frac{1}{\deg(i)}\sum_{j\sim i} w_{ij}O_{ij}v(j).$$

To encode affinity between vertices i and j (we will frequently simplify the notation  $x_i$  and  $x_j$  for vertices) consider a path  $\gamma$  of length t connecting i and j, e.g., a sequence of consecutive edges of length t, and take

$$O(\gamma) = O_{i=i_0,i_1}O_{i_1i_2}\ldots O_{i_{t-1},i_t=j}$$

This orthogonal transformation will a priori depend on the path  $\gamma$ , and we will return to this point later.

A main point of the VDM approach is the characterisation of the size of this orthogonal transformation.

Recall that we have defined a  $d \times d$  matrix

$$S_{i,j} = egin{cases} w_{ij}O_{ij}, & (i,j)\in E \ 0 & ext{otherwise} \end{cases}$$

where  $O_{ij} \in O(d)$ ,  $0 \in O(d)$ .

Affinity is a measure of compatibility between transformations obtained by following different paths from i to j.

This is measured using the matrix

$$\widetilde{S} = D^{-1/2} S D^{1/2}$$

which is similar to the matrix  $D^{-1}S$  introduced earlier.

The affinity between vertices i and j is defined as follows. For a natural number t consider all paths from i to j of length 2t, and put:

$$A_t(i,j) = \|\widetilde{S}^{2t}(i,j)\|_{HS}^2$$

where  $\| - \|$  is the Hilbert-Schmidt norm of this matrix:

$$\|T\|_{HS} = \left(\sum_{i,j} |T_{ij}|^2\right)^{1/2}$$

The matrix  $\widetilde{S}$  is symmetric, so it has a complete set of eigenvectors  $v_1, \ldots, v_{nd}$  and eigenvalues  $\lambda_1, \ldots, \lambda_{nd}$ .

Then

$$\widetilde{S}(i,j) = \sum_{l=1}^{nd} \lambda_l v_l(i) v_l(j)^T; \quad \widetilde{S}^{2t}(i,j) = \sum_{l=1}^{nd} \lambda_l^{2t} v_l(i) v_l(j)^T$$

where  $v(i) \in \mathbb{R}^d$ .

$$\begin{split} \|\widetilde{S}^{2t}(i,j)\|_{HS}^{2} &= \operatorname{Tr}(\widetilde{S}^{2t}(i,j)\widetilde{S}^{2t}(i,j)^{T}) \\ &= \sum_{l,r=1}^{nd} (\lambda_{l}\lambda_{r})^{2t} \langle v_{l}(i), v_{r}(i) \rangle \langle v_{l}(j), v_{r}(j) \rangle \\ &= \langle V_{t}(i), V_{t}(j) \rangle \end{split}$$

where

$$V_t: i \mapsto (\lambda_l \lambda_r)^{2t} \langle v_l(i), v_r(i) \rangle$$

and l, r = 1, ..., nd.

#### Cryoelectron microscopy



Picture from Picture from Singer, Shkolnitsky via Baideira *Ten lectures and forty two problems . . .* 

The pixel intensity of the projection is described as

$$P_i(x,y) = \int_{-\infty}^{\infty} \phi_i(x,y,z) dz$$

where z is the direction of the imaging electrons. Here  $\phi_i$  is defined as follows.

Let  $\phi(x, y, z)$  be the electric potential of the 'laboratory molecule', then

$$\phi_i(x, y, z) = \phi(R_i^{-1}(x, y, z)),$$

where  $R_i$  is the SO(3) matrix moving the laboratory molecule to its position when it is X-rayed.

Problem: Find the rotation matrices  $R_i$ .

Note that for two molecules i and j the element of SO(3) given by

$$\rho(i,j) = R_j R_i^{-1}$$

represents a measure of similarity between two projections.

A different version of the same question: can we work out the necessary group actions  $\rho(i, j)$ ?

#### Shape transport



Picture from Sayan Muhkerjee

# **Object** alignment



Picture from Sayan Muhkerjee

## Data Acquisition: microCT (High Resolution X-ray CT)



Surface reconstructed from  $\mu$ CT-scanned voxel data

Landmarked Teeth 
$$\rightarrow$$
  
 $d_{Procrustes}^{2} (S_{1}, S_{2}) = \min_{\substack{R \text{ rigid motion}}} \frac{1}{k} \sum_{j=1}^{k} ||R(x_{j}) - y_{j}||^{2}$ 

Boyer et al. "Algorithms to Automatically Quantify the Geometric Similarity of Anatomical Surfaces." *Proceedings of the National Academy of Sciences* 108.45 (2011): 18221-18226.

### Myriads of Shape Distances...

 $\begin{array}{ll} d_{\rm cWn}\,(S_1,S_2) & \mbox{Conformal Wasserstein Distance (CWD)} \\ d_{\rm cP}\,(S_1,S_2) & \mbox{Continuous Procrustes Distance (CPD)} \\ d_{\rm cKP}\,(S_1,S_2) & \mbox{Continuous Kantorovich-Procrustes Distance (CKPD)} \end{array}$ 

$$d_{\mathrm{cP}}\left(S_{1},S_{2}\right) = \inf_{\mathcal{C} \in \mathcal{A}\left(S_{1},S_{2}\right)} \inf_{R \in \mathbb{Z}\left(3\right)} \left(\int_{S_{1}} \left\|R\left(x\right) - \mathcal{C}\left(x\right)\right\|^{2} d\mathrm{vol}_{S_{1}}\left(x\right)\right)^{\frac{1}{2}}$$







### Shape Space



The comparisons are mostly local and work best over small distances.

# CT scans



# Geometry of synchronisation

# **Classifying objects**



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Given a sequence of objects  $(o_1, ..., o_n)$  and a group (or collection of groups) *G* the objective is to learn a collection of group elements  $\rho_{ij}$  that would transform object  $o_i$  to  $o_j$ . Synchronisation problems deal with questions regarding aligning a collection of objects in a consistent manner.

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If it is possible to find a sequence of group elements that allow for an accurate transformation between objects then this set can be *synchronised*.

#### Synchronization problem with noise

$$y_i = R_i x + \xi_i$$
  
 $R_i \in O(d), \quad \xi_i \sim \text{i.i.d. noise}$ 



Afonso S. Bandeira. "Ten Lectures and Forty-Two Open Problems in the Mathematics of Data Science." (2015). 47

# Spectral properties of graphs

- Assume given a graph G, with a set of vertices V and a set of edges E. For simplicity we assume that G has no loops or multiple edges. (This assumption can be removed.) For a vertex v we denote its degree (or valency) by d<sub>v</sub>: this is the number of edges that meet at v.
- Is it possible to measure the *level of complexity of the graph*?
- For example, how *well connected* is the graph? Does it have bottlenecks? How easy is it to disconnect it into components?

The topological structure of the graph does not capture the functional information about the system it describes. To include this information, we use edge weights. A typical example of an *edge weight* is a function  $w: V \times V \to \mathbb{R}^{\geq 0}$  such that

1. 
$$w(i,j) = w(j,i)$$
 for all  $i,j$ ;  
2.  $w(i,j) = 0$  if  $(i,j) \notin E$ ;  
3.  $w(i,i) = \sum_{j=1}^{N} w(i,j)$ .

In what follows, we will use more general weights that will be allowed to take values in groups, modules over groups, etc.
#### The discrete Laplacian

The discrete Laplace operator is defined on the space of functions  $g: V \to \mathbb{C}$  on the set of vertices V.

$$\Delta g(u) = \sum_{v \sim u} (g(v) - g(u))$$

where  $v \sim u$  means that the two vertices are adjacent.

The Laplacian is the composition of two operations:

- 'gradient' (or derivative):  $d : C(V) \rightarrow C(E)$ , df(u, v) = f(v) - f(u), where u, v are ends of one edge;
- 'divergence':  $\partial : C(E) \rightarrow C(V)$ ;

$$\partial \phi(u) = \sum_{u \sim v} \phi(u, v).$$

Since the graph is finite,  $C(V) = \mathbb{R}^{\text{number of vertices}}$ ,  $C(E) = \mathbb{R}^{\text{number of edges}}$ , and the operators d,  $\partial$ ,  $\Delta$  are given by matrices of a suitable size. Since the graph is finite,  $C(V) = \mathbb{R}^{\text{number of vertices}}$ ,  $C(E) = \mathbb{R}^{\text{number of edges}}$ , and the operators d,  $\partial$ ,  $\Delta$  are given by matrices of a suitable size.

The Laplacian matrix of  $\mathcal{G}$  is the  $N \times N$  matrix L defined as

$$[L]_{ij} = \begin{cases} d_i, & \text{if } i = j; \\ -w_{ij}, & \text{if } i \neq j \text{ and } (i,j) \in E; \\ 0, & \text{otherwise.} \end{cases}$$

It is a real symmetric matrix with non-positive entries outside the diagonal, and the sum of each column (or row) is zero.

## **Normalized Laplacian**

The normalized Laplacian is the matrix  $L_n = D^{-1/2}LD^{-1/2}$ , where D is the diagonal matrix with nonzero entries

$$d_i = \sum_{i \sim j} w(i, j).$$

## Normalized Laplacian

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$$d_i = \sum_{i \sim j} w(i, j).$$

That is,

$$[L_{\mathbf{n}}]_{ij} = \begin{cases} 1, & \text{if } i = j; \\ \frac{-w_{ij}}{\sqrt{d_i}\sqrt{d_j}}, & \text{if } i \neq j \text{ and } (i,j) \in E; \\ 0, & \text{otherwise.} \end{cases}$$

The normalized Laplacian is scale-independent, and it is more advantageous for clustering purposes.

- 1. When the graph G is connected, the kernel of the Laplacian consists of constant functions. So while  $\Delta$  is not invertible, it fails to be invertible on a one-dimensional space.
- 2. The operator  $\Delta$  is positive, so it has real, nonnegative eigenvalues  $0 = \lambda_0 < \lambda_1 \leq \cdots \leq \lambda_{n-1}$ .
- 3. The first non-zero eigenvalue  $\lambda_1$  is particularly interesting. It can measure 'the growth' of the graph; an example of this is the edge expansion.

## **Isoperimetric problems**

- A classical problem in graph theory: Remove as few edges as possible to separate the graph into two parts of almost equal size.
- There are many problems of this general type: Remove as little of the graph as possible to separate out a subset of vertices of some desired 'size'. (These are called 'separator problems'.)
- An *edge-cut* is a subset of a set of edges which disconnects the graph.
- Cheeger constants appear in isoperimetric problems involving edge-cuts. They are also important in flow or routing problems on graphs.

Let S be a subset of the vertex set of G. The edge boundary  $\partial S$  of S consists of all edges with exactly one end in S. This is the same as  $\partial \overline{S}$ , where  $\overline{S} = G \setminus S$ . What is the typical size of the boundary  $\partial S$ ? This information

indicates the level of connectedness of the graph.



The *volume* of a subgraph spanned by a subset  $S \subset V$  is:

$$\operatorname{vol}(S) = \sum_{i \in S} d_i = \sum_{i \in S} w_{ii}.$$

The *boundary* of S is the sum of the weights of the edges between vertices in S and vertices not in S:

$$\partial(S) = \sum_{i \in S, j \notin S} w_{ij}.$$

The *edge expansion* of S as the quotient

$$\phi(S) = \frac{\partial(S)}{\operatorname{vol}(S)}.$$

## **Spectral clustering**

We will be interested in the question of clustering on graphs, that is, a partition of the set of vertices V



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For this, the edge expansion is a useful indicator of quality:

$$\phi(S) = \frac{\partial(S)}{\operatorname{vol}(S)}.$$

 $\phi(S)$  is small if the vertices of S are more connected among themselves than they are to the complement of S.

## Edge expansion







#### The Cheeger constant

$$h_G = \min\left\{\frac{|\partial S|}{\operatorname{vol} S}\right\}$$

For graphs with bottlenecks, the Cheeger constant will be small. close to zero. Graphs with well connected components will have a larger Cheeger constant. Here  $|\partial S| = 2$  so  $|\partial S|/\operatorname{vol} S \to 0$  as the size of S grows.



#### We have the *isoperimetric inequality*

$$\lambda_1/2 \leq h_G < \sqrt{2\lambda_1}.$$

## Expander graphs



## Applications of spectral clustering









### Power grids as graphs



The main high voltage part (above 132 kV) of the UK power grid can be represented as a graph with 810 vertices and 1194 edges. If lower voltage nodes were included the graph would have many thousands of vertices.

#### **Functional structure**



It is crucial that the graph model captures the functional and physical structure of the grid: for example, power flow through the edges.

## The islanding problem

- Find a way to identify energetically balanced islands within a power network (preferably in real time);
- The islands should have a small overall powerflow through the boundary
- Understand the effect of removing a part of the grid on the rest of the network. This requires the use of realistic constraints in all models.



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**Step 0 (Spectral dimension)** Identify gaps in the spectrum of the Laplacian weighted by the powerflow



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#### Steps 1, 2 (Normalised Laplacian, Spectral embedding)

Use the first k eigenvectors to provide coordinates to the vertices in  $\mathbb{R}^k$ , for k = 2 and 3 to embed in  $\mathbb{R}^2$  or  $\mathbb{R}^3$ .



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#### Step 3 (Power flows)



Observe that the island with just bus 12 has the smallest boundary  $\partial(S) = 8.53$  but maximal expansion  $\phi(S) = 100\%$ , as its boundary amounts to all its (electrical) size.

68

#### Step 3 (Power flows)



Alternative power flow in the IEEE 39-bus test system and clustering into four islands.



## UK power grid



# Synchronisation

Synchronisation problems deal with questions regarding aligning a collection of objects in a consistent manner.

Given a sequence of objects  $(o_1, ..., o_n)$  and a group (or collection of groups) *G* the objective is to learn a collection of group elements  $\rho_{ij}$  that would transform object  $o_i$  to  $o_j$ .
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If it is possible to find a sequence of group elements that allow for an accurate transformation between objects then this set can be *synchronised*. Let  $\Gamma = (V, E)$  be a graph with vertex set V and edge set E. Let G be a group.

- A vertex potential is a map  $f: V \to G$ .
- An edge potential is a function  $\rho: E \to G$  which is symmetric: for every edge  $(i,j) \in E$  we have  $\rho(j,i) = \rho(i,j)^{-1}$



# **Compatible potentials**

 We say that a group potential f and an edge potential ρ are compatible across an edge (i, j) iff

 $f(i) = \rho(i,j)f(j).$ 

- We say that the two potentials ν and ρ are *compatible* iff they are compatible across every edge (i, j) ∈ E.
- It is natural to extend this definition to the case f : V → M, where M is a G-module.



Bandeira, Singer and Spielman (2013) posed the following question:

Given an edge potential  $\rho: E \to G$ , does there exits a vertex potential  $f: V \to G$  which is compatible with  $\rho$ ?

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Note that in this formulation the converse problem is easy. Given a vertex potential  $f: V \to G$ , one can define a compatible edge potential  $\rho$  by

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The problem becomes more complex when we consider functions  $f: V \rightarrow M$ , for a *G*-module *M*.

A path  $\gamma$  connecting vertices v and w is a sequence of edges

$$\gamma = (e_1, e_2, \ldots, e_n)$$

where for k = 1, ..., n,  $e_k = (i_k, j_k)$ ,  $v = i_1, j_k = i_{k+1}$  and  $w = j_n$ .



An edge potential  $\rho: E \rightarrow G$  gives rise to a map

$$\{\text{Paths in }\Gamma\} \longrightarrow G, \qquad R: \gamma \mapsto \rho(e_1) \dots \rho(e_n).$$

#### Reversal property

For e = (i, j), let  $e^{-1} = (j, i)$ , and

$$\gamma^{-1} = (e_n^{-1}, \dots, e_1^{-1}).$$

Because the edge potential  $\rho$  is symmetric:

$$R(\gamma^{-1}) = \rho(e_n^{-1}) \dots \rho(e_1^{-1})$$
  
=  $\rho(e_n)^{-1} \dots \rho(e_1)^{-1} = R(\gamma)^{-1}$ 



An edge potential  $\rho: E \to G$  gives rise to a map

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Concatenation

 $R(\gamma_1 \circ \gamma_2) = R(\gamma_1)R(\gamma_2).$ 



Let  $\omega$  be a based loop at the vertex v.

Let  $\Omega_v$  be the space of loops based at v together with the trivial loop 0 = (v, v); put R(0) = 1.

Then we have a natural map from the space  $\Omega_v$  of loops based at v to G defined by

$$\omega \mapsto R(\omega)$$



## Definition

The group element  $R(\omega)$  will be called the *holonomy* of the loop  $\omega$ . We say that the holonomy of  $\omega$  is trivial if and only if  $R(\omega) = 1$ , the identity element of G.

#### Lemma

For every vertex v,  $H(v) = R(\Omega_v)$  is a subgroup of G, called the holonomy group at v. If the graph  $\mathcal{G}$  is connected, for every two vertices v and w, the holonomy groups  $\Omega_v$  and  $\Omega_w$  are isomorphic. For a path  $\gamma$  connecting a vertex v to a vertex w, let  $V(\gamma)$  be the set of vertices along  $\gamma$ , and let  $E(\gamma)$  denote the set of edges comprising  $\gamma$ .

If  $f: V \to G$ , then  $f_{\gamma}$  is the restriction of f to  $V(\gamma)$ . Similarly,  $\rho_{\gamma}$  is the restriction of  $\rho$  to  $E(\gamma)$ .

We say that  $f_{\gamma}$  is compatible with  $\rho_{\gamma}$  along  $\gamma$  if and only if for every edge  $e \in E(\gamma)$ , e = (i, j),  $f_{\gamma}(i) = \rho(i, j)f_{\gamma}(j)$ .

It follows that

$$f_{\gamma}(\mathbf{v}) = R(\gamma)f_{\gamma}(\mathbf{w}).$$

# Synchronisation along loops

#### Lemma

Let  $R(\omega) = 1$  for a based loop  $\omega$ . Then there exists an edge potential  $f_{\omega} : V_{\omega} \to G$  along  $\omega$  which is compatible with the edge potential  $\rho_{\omega}$  along  $\omega$ .

# Synchronisation along loops

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#### Proof.

Let  $\omega = (e_1, \ldots, e_n)$ , where  $i_1 = j_n = v$ . We put  $f_{\omega}(v) = 1$ , and  $f_{\omega}(j_k) = f(i_{k+1}) = \rho(e_1) \ldots \rho(e_k)$ .

Then for every edge  $e_k = (i_k, j_k)$ ,  $f_{\omega}(i_k) = \rho(i_k, j_k)f(j_k)$ .

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Then for every edge  $e_k = (i_k, j_k)$ ,  $f_{\omega}(i_k) = \rho(i_k, j_k)f(j_k)$ .

When k = n,

$$f_{\omega}(j_n) = \rho(e_1) \dots \rho(e_n) = R(\omega) = 1$$

SO

$$f_{\omega}(j_n) = f_{\omega}(i_1) = 1$$

## Theorem (Gao-B-Mukherjee)

Let  $\mathcal{G}(V, E)$  be a connected graph and let  $\rho$  be an edge potential on  $\mathcal{G}$ . Then there exists a vertex potential  $f : V \to G$  compatible with  $\rho$  if and only if the holonomy group H(v) of some vertex v is trivial. It follows that the holonomy group of every other vertex is also trivial, as is the holonomy of every loop in  $\mathcal{G}$ . Assume that the holonomy of every loop is trivial. Choose a base vertex v and put f(v) = 1. Then for every vertex w define

$$f(w) = R(\gamma)$$

where  $\gamma$  is a path connecting v to w.



If  $\gamma_1$  is a different path connecting v to w, then  $\omega = \gamma \circ \gamma_1^{-1}$  is a loop based at v, so that

$$1 = R(\omega) = R(\gamma)R(\gamma_1)^{-1}$$

so the construction is independent of the path.



## **Basic idea**

If  $\gamma_1$  is a different path connecting v to w, then  $\omega = \gamma \circ \gamma_1^{-1}$  is a loop based at v, so that

 $1 = R(\omega) = R(\gamma)R(\gamma_1)^{-1}$ 

so the construction is independent of the path. The vertex potential f is compatible with  $\rho$  by construction.



### Corollary

Let  $\rho \in P(E, G)$  be an edge potential for which the synchronization problem has a solution, that is, there exists a vertex potential  $f \in P(V, G)$ compatible with  $\rho$ .

Then for every triangle  $(u_0, u_1, u_2)$  in  $\mathcal{G}$ 

 $\rho(u_0, u_1)\rho(u_1, u_2)\rho(u_2, u_0) = 1.$ 



In summary, the triviality of holonomy groups for a symmetric edge potential implies the following well-known cycle consistency conditions:

$$\rho(i,i) = 1 \quad \text{for all } i \in V,$$
  

$$\rho(i,j)\rho(j,i) = 1 \quad \text{for all } (i,j) \in E,$$
  

$$\rho(i,j)\rho(j,k)\rho(k,i) = 1 \quad \text{if } (i,j), (j,k), (k,i) \in E.$$

## Fibre bundles and consistency conditions



**Theorem** (Steenrod 1951, §2). If topological group G acts on Y and  $\{U_i\}, \{\rho_{ii}\}$  is a system of coordinate transformations in the space X satisfying the cycle-consistency conditions then there exists a fibre bundle  $\mathscr{B}$  with base space X, fibre Y, group G, and coordinate transforms  $\{\rho_{ii}\}$ .

## Fibre bundles and consistency conditions



Non-trivial holonomy over a based loop in the base graph is represented as a non-trivial transformation of the fibre above the base vertex.

## **Bundle representation**



# **Theorem (Gao-B-Mukherjee)** Let G be a topological group, $\mathcal{G} = (V, E)$ a connected graph, and $\rho: E \to G$ a map satisfying $\rho_{ij} = \rho_{ji}^{-1}$ for all $(i, j) \in E$ . Write $\mathfrak{U} = \{U_i \mid 1 \leq i \leq |V|\}$ for an open cover of $\mathcal{G}$ in which $U_i$ the open star of the vertex i

## **Bundle representation**



# **Theorem (Gao-B-Mukherjee)** Let G be a topological group, G = (V, E) a connected graph, and

$$\begin{split} \rho &: E \to G \text{ a map satisfying} \\ \rho_{ij} &= \rho_{ji}^{-1} \text{ for all } (i,j) \in E. \text{ Write} \\ \mathfrak{U} &= \{U_i \mid 1 \leq i \leq |V|\} \text{ for an open} \\ \text{cover of } \mathcal{G} \text{ in which } U_i \text{ the open star} \\ \text{of the vertex } i \end{split}$$

## **Bundle representation**



### Theorem (Gao-B-Mukherjee)

Then  $\rho$  defines a flat principal *G*-bundle  $\mathcal{B}_{\rho}$  over  $\mathcal{G}$  with a system of local trivializations defined on the open sets in  $\mathfrak{U}$  with constant bundle transition functions  $\rho_{ij}$  on non-empty  $U_i \cap U_j$ . Furthermore,  $\rho$ is synchronizable if and only if  $\mathcal{B}_{\rho}$  is a flat principal *G*-bundle. The fibre bundle  $\mathcal{B}_{\rho}$  associated with the connected graph  $\mathcal{G}$  and edge potential  $\rho$  is called a *synchronization bundle associated with the edge potential*  $\rho$  *over*  $\mathcal{G}$ .

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Our theorem establishes a one-to-one correspondence between equivalence classes of synchronisable potentials and flat bundles on  $\mathcal{G}$  (up to isomorphism).

### Algorithm: SynCut

Input:  $\mathcal{G} = (V, E, w), \rho \in C^1(\mathcal{G}; G)$ , number of partitions KOutput: Partitions  $\{S_1, \dots, S_K\}$ 

- 1. Solve synchronization problem over  $\mathcal{G}$  for  $\rho$ , obtain  $f \in C^0(\mathcal{G}; \mathcal{G})$
- 2. Compute  $d_{ij} = \exp\left(-w_{ij} \|f_i \rho_{ij}f_j\|\right)$  on all edges  $(i, j) \in E$
- 3. Spectral clustering on weighted graph (V, E, d) to get  $\{S_1, \cdots, S_k\}$
- 4. Solve synchronization problem within each partition  $S_j$ , "glue up" the local solutions to obtain  $f_* \in C^0(\mathcal{G}; G)$
- 5.  $f \leftarrow f_*$ , repeat from Step 2

Tingran Gao, Jacek Brodzki, Sayan Mukherjee. "The Geometry of Synchronization Problems and Learning Group

# Obstruction to synchronisation

The compatibility condition on edges is

$$f(i) = \rho(i,j)f(j)$$

To construct a useful way to measure the failure to satisfy this condition, we now consider a *G*-module *F* (equipped with an inner product), and *F*-valued functions  $f : V \to F$  on the vertex set *V*.

We want to interpret the difference

 $f(i) - \rho(i,j)f(j).$ 

# Let F be a vector space such that $G \subset GL(F)$ , and $f \in C^0(\mathcal{G}; F)$ .

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Failure of synchronisation can be measured by frustration

$$\eta(f) = \frac{1}{2} \frac{\sum_{i,j \in V} \|f_i - \rho_{ij}f_j\|^2}{\sum_{i \in V} \|f_i\|^2},$$

which can be written as a Raleigh quotient

$$\eta\left(f\right) = \frac{\langle f, L_1 f \rangle}{\langle f, f \rangle}$$

where  $L_1$  is the Graph Connection Laplacian

$$(L_1 f)_i = \frac{1}{\deg(i)} \sum_{j \sim i} (f_i - \rho_{ij} f_j)$$

Given a graph  $\mathcal{G}(V, E)$  we define the gradient  $d : \mathbb{C}[V] \to \mathbb{C}[E]$  by the formula

$$df(i,j) = f(i) - f(j).$$

The divergence operator  $\delta : \mathbb{C}[E] \to \mathbb{C}[V]$  is given by

$$\delta\phi(i) = \sum_{j\sim i} \phi(i,j).$$

The graph Laplacian is the operator  $\Delta : \mathbb{C}[V] \to \mathbb{C}[V]$  defined by  $\Delta_{\rho} = \delta d$ :

$$\Delta f(i) = \sum_{j:j \sim i} (f(i) - f(j))$$

M is a G-module equipped with an inner product, and  $\rho$  a symmetric potential as before.

$$d_{\rho}: \mathbb{C}(V, M) \longrightarrow \mathbb{C}(E, M)$$
$$d_{\rho}f(i, j) = f(i) - \rho(i, j)f(j)$$
$$\in \mathbb{C}(V, M).$$

for every  $f \in \mathbb{C}(V, M)$ .
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ho}:\mathbb{C}(V,M)\longrightarrow\mathbb{C}(E,M)$$

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for every  $f \in \mathbb{C}(V, M)$ .

A solution to the synchronisation problem is an element of the kernel of  $d_{\rho}$ .

The usual divergence:

$$\delta : \mathbb{C}(E, M) \longrightarrow \mathbb{C}(V, M)$$
  
 $\delta \phi(i) = \sum_{j:j \sim i} \phi(i, j)$ 

together with  $d_{
ho}$  give a twisted Laplacian

$$\Delta_{\rho} = \delta d_{\rho} = \sum_{j:j\sim i} (f(i) - \rho(i,j)f(j))$$

$$\eta_{\rho}(f) = \frac{\langle f, \Delta_{\rho} f \rangle}{\langle f, f \rangle} = \frac{\sum_{(i,j) \in E} w_{ij} \|f(i) - \rho(i,j)f(j)\|^2}{\sum_{i \in V} d_i \|f(i)\|^2}$$

 $\eta_{\rho}(f)$  is defined in Bandeira et al. (2013) as the *frustration* 

We have defined a simple de Rham-type complex

$$\mathcal{C}: 0 \longrightarrow \mathbb{C}(V, M) \xrightarrow{d_{\rho}} \mathbb{C}_{\rho}(E, M) \longrightarrow 0.$$

The cohomology of this complex is

$$H^0(\mathcal{C},d) = \ker d_\rho$$

in degree zero. In degree one, we have that

$$H^1(\mathcal{C}, d) = \mathbb{C}(E, M)/d_{\rho}(\mathbb{C}(V, M)).$$

Thus solutions to the synchronisation problem form the zeroth cohomology group, which is the same as the kernel of  $d_{\rho}$ .

### Theorem

1. The space of solutions of the synchronisation problem given by a unitary edge potential  $\rho$  is isomorphic to the space of harmonic functions f, i.e., functions f with the property  $\Delta_{\rho}f = 0$ . Moreover, we have the following orthogonal decomposition:

$$\mathbb{C}(V,M) = \ker d_{\rho} \oplus \operatorname{Im} \delta.$$

2. The Laplace operator  $\Delta_{\rho}$  is self-adjoint and positive.

The two linear spaces  $\Omega^0(\mathcal{G}, M)$  and  $\Omega^1_\rho(\mathcal{G}, M)$  are equipped with inner products naturally induced from the *G*-invariant inner product on *M*:

$$\langle f, g \rangle = \sum_{i \in V} d_i \langle f_i, g_i \rangle, \quad \forall f, g \in \Omega^0 (\mathcal{G}, M),$$
  
 $\langle \phi, \psi \rangle = \sum_{(i,j) \in E} w_{ij} \langle \phi_{ij}, \psi_{ij} \rangle, \quad \forall \phi, \psi \in \Omega^1 (\mathcal{G}, M),$ 

With respect to these inner products, the twisted differential  $d_{\rho}$  and the divergence  $\delta$  are adjoints of each other.

Given a group G acting on a set X, simultaneously learn a new action of G on X and a partition of X into disjoint subsets  $X_1, \dots, X_K$ , such that the new action is as close as possible to the given action and cycle-consistent on each  $X_i$   $(1 \le i \le K)$ .

## **Discrete Hodge theory**

 $\mathcal{G} = (V, E)$ :

Spaces of cochains:

$$egin{aligned} \Omega^0\left(\mathcal{G}
ight) &:= \left\{f: V o \mathbb{K}
ight\}, \ \Omega^1\left(\mathcal{G}
ight) &:= \left\{\omega: E o \mathbb{K} \mid \omega_{ij} = -\omega_{ji} \; orall i \sim j
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A cochain complex

$$0 \rightleftharpoons \Omega^{0}(\mathcal{G}) \xleftarrow[\delta]{d}{\longrightarrow} \Omega^{1}(\mathcal{G}) \rightleftharpoons 0,$$

where

$$egin{aligned} \left(df
ight)_{ij} &= f_i - f_j, \quad orall f \in \Omega^0\left(\mathcal{G}
ight), \ \left(\delta\omega
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## **Discrete Hodge theory**

The standard cochain complex

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ight). \end{aligned}$$

The graph Laplacian:

$$egin{aligned} &(L_0f)_i := (\delta df)_i \ &= rac{1}{\deg{(i)}} \sum_{j \sim i} \left(f_i - f_j
ight) \quad orall i \in V, \ orall f \in \Omega^0\left(\mathcal{G}
ight). \end{aligned}$$

## Twisted De Rham-Hodge Theory

$$(L_0 f)_i = \frac{1}{\deg(i)} \sum_{j \sim i} (f_i - f_j), \quad \forall f : V \to \mathbb{K}$$
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Naïvely:

$$egin{aligned} & \left(d_{
ho}f
ight)_{ij} &= f_i - 
ho_{ij}f_j, \quad orall f \in C^0\left(\mathcal{G};F
ight) \ & \left(\delta_{
ho}\omega
ight)_i = rac{1}{\deg\left(i
ight)}\sum_{j\sim i}\omega_{ij}, \quad orall \omega \in C^1\left(\mathcal{G};F
ight) \end{aligned}$$

then  $L_1 = \delta_{\rho} d_{\rho}$ .

We have defined:

$$egin{aligned} & (d_{
ho}f)_{ij} \sim f_i - 
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**There is a problem**:  $d_{\rho}$  does not map into  $C^{1}(\mathcal{G}; F)$  (no skew-symmetry).

$$f_j - \rho_{ji}f_j = -\frac{\rho_{ji}}{\rho_{ji}}(f_i - \rho_{ij}f_j) \neq -(f_i - \rho_{ij}f_j).$$

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**Fix**: Interpret  $f_i - \rho_{ij}f_j$  as the "local expression" of  $(d_\rho f)_{ij}$  in a local trivialization over  $\mathfrak{U} = \{U_i \mid 1 \le i \le |V|\}$  of the associated *F*-bundle of  $\mathscr{B}_{\rho}$ , denoted as  $\mathscr{B}_{\rho}[F]$ , such that the extra  $\rho_{ji}$  factor encodes a bundle transformation from  $U_i$  to  $U_i$ .

• Combinatorial Hodge Theory:

$$0 \rightleftarrows \Omega^{0}\left(\mathcal{G}\right) \xleftarrow{d}{\delta} \Omega^{1}\left(\mathcal{G}\right) \rightleftarrows 0,$$

• *Twisted* Combinatorial Hodge Theory:

$$0 \rightleftharpoons C^{0}(\mathcal{G}; F) \xleftarrow{d_{\rho}}{\longleftrightarrow} \Omega^{1}(\mathcal{G}; \mathscr{B}_{\rho}[F]) \rightleftharpoons 0.$$

# **Theorem (Gao, B, Mukherjee (2016)))** Define

$$\Delta^{(0)}_{
ho} := \delta_{
ho} d_{
ho}, \quad \Delta^{(1)}_{
ho} := d_{
ho} \delta_{
ho}$$

then the following Hodge-type decomposition holds:

$$C^{0}(\mathcal{G}; F) = \ker \Delta_{\rho}^{(0)} \oplus \operatorname{Im} \delta_{\rho} = \ker d_{\rho} \oplus \operatorname{Im} \delta_{\rho},$$
$$\Omega^{1}(\mathcal{G}; \mathscr{B}_{\rho}[F]) = \operatorname{Im} d_{\rho} \oplus \ker \Delta_{\rho}^{(1)} = \operatorname{Im} d_{\rho} \oplus \ker \delta_{\rho}.$$

Solutions to the synchronisation problem are elements of ker  $d_{\rho}$ .

Learning group actions

Given a group G acting on a set X, simultaneously learn a new action of G on X and a partition of X into disjoint subsets  $X_1, \dots, X_K$ , such that

- the new action is as close as possible to the given action;
- and is cycle-consistent on each  $X_i$   $(1 \le i \le K)$ .

If the set X is a vector space and we seek a direct sum decomposition  $X = \bigoplus_{i=1}^{K} X_i$ , the LGA problem reduces to the search for all irreducible *G*-subrepresentations of *X*.

$$X = \{x_1, \dots, x_n\}$$
 equipped with  
 $S : X \to \{\pm 1\}.$   
Let  $G = \{\pm 1\}$  act on X transitively as  
 $(g_{ji}, x_i) \mapsto x_j, g_{ji} = S(x_j)S(x_i).$ 



Suppose the spin of each point in X (i.e. the label map S) is unknown, but we know the group actions  $\{g_{ij}\}$ Then we can reconstruct S — by spectral clustering the dataset X, viewed as vertices of a complete graph  $\mathcal{G}$  with weight  $w_{ij} = g_{ij}$  on the edge connecting  $x_i$  and  $x_j$ .



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## Lemurs



#### Lemurs



If we restrict attention to the determinant of the transformation providing the best fit then we discover that teeth 1, 2, 7, 8 belong to one cluster (right), while 3, 4, 5, 6 belong to the other cluster (left).

## Identifying primate genera



## Classification into folivorous, frugivorous, and insectivorous



### **Extensions: defects in crystals**



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## 



The local structure of the defects is encoded in a comparison map  $\phi : \mathbb{R}^2 \to \mathbb{R}^2$  which compares an ideal crystal to one with defects.



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action of point groups which determine symmetries of the crystal.

The geometry of a planar crystal is determined by a a two-dimensional de Rham complex.



# The shape of lungs in COPD
The shape of lungs in chronic obstructive pulmonary disease, 2017 F. Belchi Guillamon, M. Pirashvili, M. Bennett, J. Conway, R. Djukanović, J. Brodzki Scientific Reports, 28 March 2018, supported by the EPSRC Joining the Dots programme



• Chronic obstructive pulmonary disease (COPD) is a common progressive disease, affecting more than 200 million individuals worldwide and is the third leading cause of death.

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- Cigarette smoking is the main COPD risk factor, but it is quite frequent among people who never smoked, and many smokers do not develop COPD.
- Many adults who develop COPD have low lung function in the early life.

# CT scans



#### Related work: Arterial trees in the brain



Bendich et al., Ann. Appl. Statistics, 2016

## The geometry of the lung



Pictures from Smith et al., Human airway branch variation and chronic obstructive pulmonary disease, PNAS 2018

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#### Definition

Let X and Y be finite multisets, whose underlying sets are subsets of a metric space Z. Assume that X and Y are of equal size as multisets.

The bottleneck distance between X and Y is defined by

$$d_B(X,Y) = \inf_{\gamma} \sup_{x \in X} d(x,\gamma(x))$$

where  $\gamma$  ranges over all bijections of multisets  $\gamma: X \to Y$ .

We consider each point of multiplicity k as k points, and  $\gamma$  is a bijection between the resulting sets.

#### The bottleneck distance



White and black dots belong to persistence diagrams of two different functions. The squares represent a bijection realising the bottleneck distance, which equals half the side length. Let X, Y be subsets of  $\mathbb{R}^2$  (multisets in general).

Let  $\gamma: X \to Y$  be a bijection. For  $p \ge 1$  the cost of  $\gamma$  is defined by

$$C_{p}(\gamma) = \left(\sum_{x \in D(f)} \| x - \gamma(x) \|^{p}\right)^{1/p}$$

Definition

$$W_p(X, Y) = \inf_{\gamma} C_p(\gamma)$$

where the infimum is taken over all bijections  $\gamma: X \to Y$ .

- The overall aim of this study was to develop a set of new radiomic features that can distinguish between healthy non-smokers as well as healthy smokers and patients with COPD.
- The following four study participant groups defined by smoking status and spirometry given by the GOLD guidelines were studied:
  - healthy non-smokers and healthy smokers (both judged as healthy by spirometry showing FEV1 > 80% of predicted and FEV1/FVC > 0.75),
  - mild COPD patients, consisting of GOLD stage 1 (with FEV1  $\geq$  80% of predicted and FEV1/ FVC < 0.70)
  - moderate COPD patients, consisting of GOLD stage 2 (50%  $\leq$  FEV1 < 80% of predicted and FEV1/FVC < 0.70).







Directional complexity: the sum of the lengths of the bars in the graph (total life span of the persistence diagram).



Directional complexity distinguishes healthy subjects from Mild and Moderate COPD

#### Comparison to other standard measurements



#### Comparison to other standard measurements









Numerical invariant: total barcode length in the  $H_1$  barcode.



Metric space structure created using the Wasserstein metric and represented through 2-dim MDS projection









Metric space structure created using the Wasserstein metric; represented here in a 2-dim MDS projection







- Persistence for fibre bundles: a combination of geometric and topological information
- Synchronisation through persistence
- Symmetry detection
- Time evolution: geometric intuition to guide the evolution of persistent homology
- Phase transitions in persistence diagrams

The same equations have the same solutions. Richard Feynman